

14 October 2021

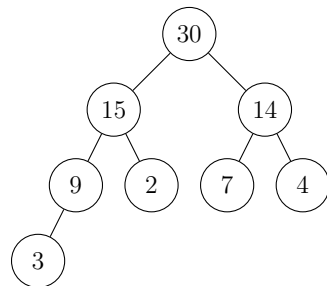
This worksheet uses the following definitions.

- **height of a node:** length of the longest path from a node to a leaf below it
- **height-balanced (AVL) tree:** satisfies $|\text{height}(\text{right child}) - \text{height}(\text{left child})| \leq 1$

1. **Warm up:** Answer the following questions about tree structures.

- What is the shortest and longest path between two level ℓ nodes in a binary tree?
- What is the smallest and largest numbers of leaves a height h binary tree can have?
- What is the heap property of a tree?
- What is the advantage of using heaps for sorting?

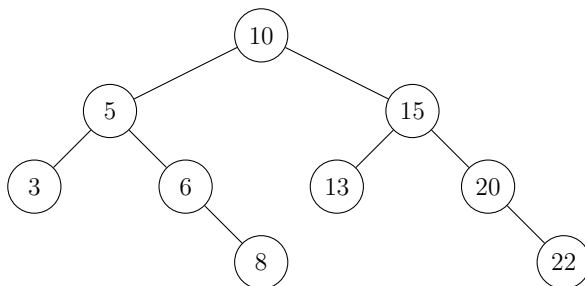
2. In the **max-heap** below left, perform the operations below right, in the given order.



- insert 16
- insert 17
- remove max
- insert 18
- insert 19

(f) Give the 0-based array that stores the heap after the steps above have been applied.

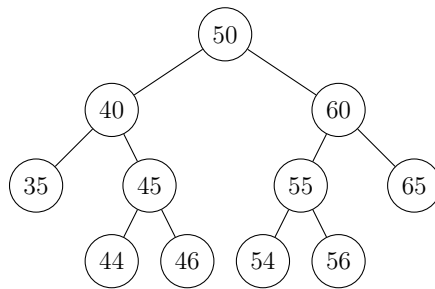
3. In the **binary search tree** below left, perform the operations below right, in the given order. For each insert operation, show the path taken when searching for the spot to insert the key.



- insert 7
- insert 14
- delete 15
- insert 15
- delete 10
- insert 10

(g) Give the DFS postorder of the tree after the steps above have been applied.

4. In the **AVL tree** below left, perform the operations below right, in the given order. Make sure to rebalance (if necessary) after every operation.



- (a) insert 43
- (b) insert 66
- (c) insert 67
- (d) remove 60

Recall that rebalancing is done in terms of **rotations**. The *x-over-y* rotation of T , for nodes x, y of T where x is a child of y , is a new tree T' identical to T , except for:

- if $T.y.parent = z$, then $T'.x.parent = z$ and $T'.y.parent = x$
- if $T.y.leftchild = x$, then $T'.y.leftchild = T.x.rightchild$
- if $T.y.rightchild = x$, then $T'.y.rightchild = T.x.leftchild$